

## quiz 09.1 Solution

(1) Using the alphabet  $GF(13)$  find a generator matrix for a Reed-Solomon code correcting any 5 bit errors. (Use primitive root 2 mod 13.)

To correct any  $e$  bit errors, the minimum distance must be at least  $2 \cdot e + 1$ . Thus, to correct any 5 bit errors, we should take designed distance  $2 \cdot 5 + 1 = 11$ . Thus, we take generating polynomial

$$g(x) = (x - 2)(x - 2^2)(x - 2^3)(x - 2^4) \dots (x - 2^{11-2})(x - 2^{11-1})$$

The length of any RS code using alphabet  $GF(13)$  is 13-1, so this code will have length 12. Multiplying together these linear factors gives

$$\begin{aligned} &g(x) \\ &= x^{10} + 8x^9 + 5x^8 + 10x^7 + 6x^6 + 4x^5 + 3x^4 + 9x^3 + 12x^2 + 7x + 11 \end{aligned}$$

Inserting the coefficients in ascending order into a 2-by-12 matrix, padding at the right with 0's in the first line and cycling over until the constant coefficient bumps against the right edge, gives the generating matrix

$$\begin{pmatrix} 11 & 7 & 12 & 9 & 3 & 4 & 6 & 10 & 5 & 8 & 1 & 0 \\ 0 & 11 & 7 & 12 & 9 & 3 & 4 & 6 & 10 & 5 & 8 & 1 \end{pmatrix}$$

Alternatively, one might also try a little trick to reduce the computational load. Since

$$x^{12} - 1 = (x - 2^1)(x - 2^2) \dots (x - 2^{13-1})$$

The check polynomial  $h(x)$  is the coefficient-reversed form of  $(x^{12} - 1)/g(x)$ . Using the factored expression saves work:  $h(x)$  is simply (the coefficients-reversed form of) the leftover factors from  $x^{12} - 1$  after the factors from  $g(x)$  are removed:

$$h(x) = (\text{reversed...})(x - 2^{11})(x - 2^{13-1}) = (\text{reversed...})x^2 + 5x + 7 = (\text{reversed...})1 + 5x + 7x^2$$

Thus, the generator polynomial is perhaps best computed by dividing  $x^{12} - 1$  by (re-reversed)  $h(x)$ , obtaining

$$\begin{aligned} g(x) &= \frac{x^{12} - 1}{x^2 + 5x + 7} \\ &= x^{10} + 8x^9 + 5x^8 + 10x^7 + 6x^6 + 4x^5 + 3x^4 + 9x^3 + 12x^2 + 7x + 11 \end{aligned}$$

It is not clear which is easier, the division, or simply multiplying out.

(2) Show that the binary fifth-degree polynomial 101001 (with coefficients in descending order) is irreducible.

We do this by trial division. That is, we know that if there is a divisor  $d(x)$  of 101001 with

$$0 < \deg(d) < \deg(101001) = 5$$

then in fact there is a divisor with

$$0 < \deg(d) \leq \deg(101001)/2$$

Thus, we need only consider possible divisors of degrees 1 and 2, since degrees are integers. Further, we need only consider irreducible divisor candidates. There are two polynomials of degree 1,  $x$  and  $x + 1$ . From class (or another family of trial divisions) among the 4 quadratic polynomials there is only one irreducible,  $x^2 + x + 1$ . Thus, we have only 3 trial divisions to do to verify irreducibility. Division by  $x$  gives remainder 1, as does division by  $x + 1$ . Division by  $x^2 + x + 1$  gives  $x + 1$ , which is not 0. Thus, 101001 is indeed irreducible.

(3) Find an element of order 63 in  $GF(2^6)$ , where  $GF(2^6)$  is modeled as  $\mathbf{F}_2[x] \text{ mod } 1010111$  where those are the coefficients in descending order.

Let  $P(x) = 1010111$  construed as the coefficients in descending order. Implicit in the question is the apparent fact that  $P(x)$  is irreducible, or else polynomials modulo  $P(x)$  wouldn't be a field. In light of the facts mentioned in the notes, namely the analogue of Fermat's little theorem for finite fields, we know that for any polynomial  $g$  not divisible by  $P$ ,  $g^{63} = 1 \pmod{P}$ . So to check whether or not a given candidate  $g$  is a primitive element, we need to compute  $g^{63/3}$  and possibly  $g^{63/7} \pmod{P}$  since 3,7 are the two primes dividing 63. In this example, the polynomial  $g(x) = x$  is not primitive, because  $x^{21} = 1 \pmod{P}$ . But it does turn out that  $x + 1$  is primitive, because neither  $(x + 1)^{21}$  nor  $(x + 1)^9$  is 1 mod  $P$ . The fast exponentiation computations to compute  $(x + 1)^9$  and  $(x + 1)^{21} \pmod{1010111}$  are

$$\begin{array}{rcl} x + 1 & 9 & 1 \\ x + 1 & 8 & x + 1 \\ x^2 + 1 & 4 & x + 1 \\ x^4 + 1 & 2 & x + 1 \\ x^3 + x & 1 & x + 1 \\ x^3 + x & 0 & x^4 + x^3 + x^2 + x \end{array}$$

$$\begin{array}{rcl} x + 1 & 21 & 1 \\ x + 1 & 20 & x + 1 \\ x^2 + 1 & 10 & x + 1 \\ x^4 + 1 & 5 & x + 1 \\ x^4 + 1 & 4 & x^5 + x^4 + x + 1 \\ x^3 + x & 2 & x^5 + x^4 + x + 1 \\ x^4 + x + 1 & 1 & x^5 + x^4 + x + 1 \\ x^4 + x + 1 & 0 & x^5 + x^3 + x^2 + x + 1 \end{array}$$

showing that  $x + 1$  to these powers is not 1, so it is a primitive element.

(4) Model the finite field  $\mathbf{F}_{32} = GF(32)$  as  $\mathbf{F}_2[x]/f(x)$ , where  $f(x)$  is the polynomial with coefficients 101001 in descending order. Find the multiplicative inverse of 1011, again interpreted as a polynomial with coefficients in descending order.

Use the extended Euclidean algorithm applied to polynomials  $g(x) = 1011$  and  $f(x) = 101001$  to obtain polynomials  $a(x)$  and  $b(x)$  so that  $a(x)f(x) + b(x)g(x) = 1$ . Looking at the latter relation modulo  $f(x)$  gives  $b(x)$  as a multiplicative inverse of  $g(x)$  modulo  $f(x)$ . Executing the extended Euclidean algorithm:

$$(x^5 + x^3 + 1) + (x^2)(x^3 + x + 1) = x^2 + 1$$

$$(x^3 + x + 1) + (x)(x^2 + 1) = 1$$

$$(x^2 + 1) + (x^2 + 1)(1) = 0$$

$$\begin{aligned} 1 &= (1)(x^3 + x + 1) + (x)(x^2 + 1) \\ &= (1)(x^3 + x + 1) + (x)[(x^5 + x^3 + 1) + (x^2)(x^3 + x + 1)] \\ &= (x)(x^5 + x^3 + 1) + (x^3 + 1)(x^3 + x + 1) \end{aligned}$$

Thus, by the earlier comment about expressions of the form  $as + bt = 1$ , the multiplicative inverse of 1011 modulo 101001 is  $x^3 + 1$ .