Math 4707 Oct. 7,2020  
Note: There are two resources in PDF  
on generating functions on the syllabus:  
One by G. Musiker (mtro)  
S. Hoptons (exercises)  
Fibonacci numbers  
recurrente Fh+1 = Fn+ Fn-1 for n21  
initial 
$$F_0 = 0$$
  $\frac{m/Fn}{2}$   
initial  $F_1 = 1$   $\frac{m/Fn}{2}$   
Noticed that  $\frac{1}{2}$   $\frac{3}{2}$   
 $F_3$  is even  $\frac{3}{2}$   $\frac{1}{3}$   
 $F_3$   $F_5$  is divisible by S  $\frac{1}{10}$   $\frac{5}{55}$   
 $F_5$  Guessed: Fe Fkm  $\forall$  m21  
 $k_{21}$ 

THEOREM: 
$$F_{k}$$
 [Fkn for k, m ≥1  
poot: We deduce it from  
LEMMA: Factor = Factor For + Fato  
for a, b ≥0  
via choosing a:= k(m-1)  
b:= k-1  
then Lemma says  
Factor 1 = Factor For + Fato  
II = Factor + Fato  
Factor 1 = Factor For + Fato  
Factor 1 = Factor For + Fato  
II = Factor + Fato  
II = Fato

It remains to prove the LEMMA.

UPMA: Farber = Far Fort + Fats  
Left's check the motances  

$$a=0$$
:  $F_{ber} = F_1 \cdot F_{ber} + F_0 \cdot F_b$   
 $= F_{ber} \cdot says ustring!$   
 $a=1$ :  $F_{ber} = F_2 \cdot F_{ber} + F_1 \cdot F_b$   
 $F_{ber} = F_2 \cdot F_{ber} + F_0 \cdot F_b$  is the original  
 $a=2$ :  $F_{ber3} = F_3 \cdot F_{ber} + F_2 \cdot F_b$   
 $F_{ber3} = 2 \cdot F_{ber} + F_b \cdot seems new$   
 $F_{ber3} \cdot F_b \cdot F_{ber}$   
 $F_{ber3} \cdot F_{ber} = f_{ber}$   
 $F_{ber3} \cdot F_b \cdot F_{ber}$ 

§4.3 An exact formula for Fn  
Weill get this formula by constraining  
a power series  

$$f(x) = F_0 \cdot x^\circ + F_1 \cdot x^\circ + F_2 \cdot x^2 + F_3 \cdot x^3 + ...$$
  
 $= \sum_{n=0}^{\infty} F_n x^n$   
 $= 0 \cdot x^\circ + 1 \cdot x^1 + 1 \cdot x^2 + 2 \cdot x^2 + 3 \cdot x^3 + 5 \cdot x^5 + 5 \cdot x^6 + ...$   
called the generating function for  
the sequence  $\{F_n\}_{n=0,1,2,-..}$   
e.g.  $a_n = 2^n = \#$  subsets of  $\{1,2,-n\}$   
has generating function  
 $a(x) = 2^{\circ} \cdot x^\circ + 2^{1 \cdot x^1} + 2^{\circ} \cdot x^2 + 2^{\circ} \cdot x^2 + ...$   
 $= 1 + 2x + 9x^2 + 8x^3 + ...$ 

$$= \frac{a}{1-r}$$

e.g. Fix n, and consider  

$$I = b_{k} = \binom{n}{k}$$

$$I \ge I = \binom{n}{2} \binom{n}$$

**-** -

THEOREM: Fibonacci #'s 
$$\{F_n\}$$
 have  
(a)  $f(x) = \sum_{n=0}^{\infty} F_n x^n = 0 \cdot x^n t \cdot 1 \cdot x^1 + 1 \cdot x^2 + 2x^3$   
 $= \frac{x}{1 - x - x^2}$   
(b) ... and from these we will deduce  
 $F_n = \frac{1}{\sqrt{5}} (\phi^n - \beta^n)$  where  
 $\phi = \frac{1 \pm \sqrt{5}}{1 \pm \sqrt{5}} \approx 1.618... = \operatorname{golden} > 1$   
 $\beta = \frac{1 \pm \sqrt{5}}{2} \approx -0.618..., (\beta) < 1$   
(c) ... and hence  $F_n \approx \frac{1}{\sqrt{5}} \phi^n$  as  $n \to \infty$   
and fins  $F_{1n} = \phi = g^{nden neto}$ .  
 $n \to \infty$   $F_n = \left[\sqrt{5} \phi^n\right] \quad \forall n \ge 0$   
 $\operatorname{fins} F_{n} = \left[\sqrt{5} \phi^n\right] \quad \forall n \ge 0$   
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Asit DE  

$$p = \frac{1+15}{2} = \frac{1}{2} \frac{$$

THEOREM: Fibonacci #'s 
$$\{F_n\}$$
 have  
(a)  $f(x) = \sum_{n=0}^{\infty} F_n x^n = 0 \cdot x^n + 1 \cdot x^1 + 1 \cdot x^2 + 2x^3 + 5x^4 + 1 \cdot x^2 + 2x^3 + 1 \cdot x^3 + 1 \cdot x^4 + 1 \cdot$ 

$$= \frac{\chi}{1-\chi-\chi^2}$$



(b) 
$$f(x) = \frac{x}{1-x-x^2}$$
  
from these we will deduce  
 $F_n = \frac{1}{\sqrt{5'}} (\phi^n - \beta^n)$  where  
 $\phi = \frac{1}{\sqrt{5'}} (\phi^n - \beta^n)$  where  
 $f(x) = \frac{x}{1-x-x^2} = (1 - \phi x)(1 - \beta x)$   
 $(1 - (1 + \sqrt{5} + \frac{1}{\sqrt{5}})x + \frac{1}{\sqrt{5}} + \frac{1}{\sqrt{5}})x + \frac{1}{\sqrt{5}} + \frac{1}{\sqrt{5}}$ 

(c) 
$$F_n = \frac{1}{\sqrt{5}} \left( p^n - p^n \right)$$
  $p = \frac{1}{2} \frac{1}{5} \frac$ 

$$S(3,1) = 1 = \#\{123\}$$
  

$$S(3,2) = 3 = \#\{12-3, 13-2, 23-1\}$$
  

$$S(3,3) = 1 = \#\{1-2-3\}$$

$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	n	1	2	3	4	5	
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	J	1					
3 1 3 1	2	1	1				
	3	1	3	1			
4 1 7 6 1	4	1	7	6	1		-
5 1 15 25 10 1	5	1	15	25	10	1	

#2(a) 
$$S(n,1)=1=S(n,n)$$
  
only all of  
 $l_{1,2,-5n}$  only  $1-2-3---n$   
 $early massingle block block block
(b)  $S(n,n-1)=\binom{n}{2}$   
 $(b) S(n,n-1)=\binom{n}{2}$   
 $(b) S(n,n-1)=\binom{n}{2}$   
 $(c) S(n,2)=\frac{2^{N}-2^{C}}{32}$   $(=2^{n-1}-1)$   
 $(numerator picks a)$   
 $subset S \neq 0, (1,2,-n)$   
 $b be the 1st block bloc$$ 

#3. 
$$S(n_1k) = S(n-1, k-1) + k S(n-1, k)$$
  
in goes into a  
ingleton block  
 $11/2 - 2^{n-1}$   
ingleton block  
 $11/2 - 2^{n-1}$   
induces  
 $11/2 - 2^{n-1}$   
induces  
 $11/2 - 2^{n-1}$   
induces  
 $11/2 - 2^{n-1}$   
induces  
 $11/2 - 2^{n-1}$   
 $1/2 - 2^{n-1}$   
 $1/2 + 2^{$ 

$$\begin{array}{l} \#5. \quad S(n_{1}k_{2}) = S(n_{-1}k_{-1}) + kS(n_{-1}k_{2}) \quad \text{for} \\ & n \geq k \geq q \\ \sum S(n_{1}k_{2}) \times = \sum S(n_{-1}k_{-1})^{n} + \sum kS(n_{-1}k_{2})^{n} \\ & n \geq k \\ n \geq k \\ \end{array}$$

$$\begin{array}{l} f_{k}(x) = \chi f_{k-1}(x) + \chi \sum kS(n_{-1}k_{2})^{n-1} \\ & n \geq k \\ \end{array}$$

$$\begin{array}{l} f_{k}(x) = \chi f_{k-1}(x) + \chi \times f_{k-1}(x) \\ & f_{k}(x) = \chi f_{k-1}(x) + \chi \times f_{k-1}(x) \\ & (1 - k_{x})f_{k}(x) = \chi f_{k-1}(x) \\ & f_{k}(x) = \frac{\chi}{1 - k_{x}} \\ & = \frac{\chi}{1 - k_{x}} \cdot \frac{\chi}{(-k_{-1})_{x}} f_{k-2}(x) \\ & = \frac{\chi}{1 - k_{x}} \cdot \frac{\chi}{(-k_{-1})_{x} - \frac{\chi}{1 - \chi}} \\ & = \frac{\chi}{(1 - k_{x})(1 - (k_{-1})_{x}) - (1 - 2x)(1 - x)} \end{array}$$

e.g.  

$$f_{4}(x) = S(t_{4}, u) x^{4} + S(5, u) x^{5} + S(6, u) x^{b} + ...$$

$$= \sum S(n, 4) x^{n}$$

$$= \frac{X^{4}}{(1 - x)(1 - 2x)(1 - 3x)(1 - 4x)}$$
functions  

$$= \frac{A}{1 - x} + \frac{B}{1 - 2x} + \frac{C}{1 - 3x} + \frac{D}{1 - 4x}$$

$$= A \sum_{v \ge 0} x^{v} + B \sum_{v \ge 0} 2^{v} x^{v} + C \sum_{v \ge 0} 3^{v} x^{v} + D \sum_{v \ge 0} 4^{v} x^{v}$$

$$= \sum (A + 2^{v}B + 3^{v}C + 4^{v}D) x^{v}$$

$$= S(u, 4)$$
has a formula under u

MATH 4707 Oct. 14, 20 GRAPH THEORY Enler walks and Enler tours (\$7.3) EXAMPLE: The 7 bridges of Königsberg land wass Kneiphot and mass land mass Q: Can one find a tour (= a closed walk) Some starting and ending land mass that goes over each bridge exactly once? Can one find a walk (possibly different start/end masses) using each bridge exactly once?

land mass and mass neiphot land mass 7 L. Enler proved both are impossible in 1736, and came up with a graph theory abstraction that solved all such publicus! vertices edges V.E) A graph G= Abobraction: (all!) (bridges) vertex { and masses dear note that we're lana multiple ec with same endports! DEFINITION: An Enter tour in Gi is a closed walk from vertex-to-vertex not simple graphs along edges, using each edge exactly once. An Enter walk/path is some but not necessarily

Q: Which graphs have Enlertours, Euler walks? Are they unique in any sense? all even dega(u) Yes, has an Euler tour. mique 10 degree vertices odd No Euler tours, but has Euler walks. THEORET (Fuler 1736) G= (V, G) a graph with no isolated vertices has an Ruler tour (a) G is connected i.e. every pair v, v'∈V AND has some path v between them (b) every vertex ver has deg G (v) even degree .

THEORET (Fuler 1736) G= (V, E) a grouph with no isolated vertices has an Euler town (a) G is connected i.e. every pair v, u'EV AND has some path V between them (b) every vertex ver has deg G (v) ever degree of v proof: (=>): Close off the Enter tour, and direct it with arrows: Gran any vivier pick edges incident to theme, e' and following the intimite tour from e to e' gives a parts from utor. So G is connected. Also, fixing vel, the four pours off edges four pours off edges incident to v as it enters v and exits v along those edges.

THEOREM (Enler 1736) G= (V, E) a grouph with no isolated vertices has an Euler town v, v EV (a) G is connected i.e. every pair v, v EV has some path v Letween them (b) every vertex ver has deg G (v) even proof: (=): Given a graph & that is connected and has all degg(u) even, let's give an algorithm to find an Erler tour. Start at some vertex vo EV, more along anedge to some V, and erase the edge you used. Then go from V, to some v2 along an unnsed edge. Repeat until you get stuck at some vertex vk until you get stuck at some vertex vk uhere all its incident edges were eroused. CLAIM: VE=Vo Only vo has odd degree in the erased edge gruph, as all others maintan even degree as you enter und Leave.

